

Overview

- In-depth look at vacuum in Postgres 9.4
- Code references in slide notes

Topics

- MVCC
- CLOG and MultiXacts
- What can be vacuumed
- Freezing
- HOT
- Vacuum
- Autovacuum

MVCC

- Postgres uses Multi Version Concurrency Control
- Rows are never deleted, they are only marked as deleted as of a specific transaction
- Updates are essentially a delete of the old tuple and an insert of a new tuple with the new values
- Eventually these old, dead tuples must be removed by "vacuuming"

CLOG & MultiXacts

- The CommitLOG tracks transaction status (committed, aborted, in-progress)
- MultiXacts store information about row-level locks, updates and deletes when multiple transactions are involved.
- Vacuum is responsible for removing unnecessary data from CLOG and MultiXacts

CLOG and MultiXacts are stored as SLRUs (Simple Least Recently Used).

See src/backend/access/transam/clog.c, multixact.c and slru.c

What can be vacuumed?

- Vacuum can only remove rows that no currently running transaction could see.
- Generally limited by **oldest running transaction** in the database
- Might also be limited by streaming replication (with *hot_standby_feedback* enabled), prepared transactions, or logical decoding
- Some special handling for current XIDs and locks

See HeapTupleSatisfiesVacuum() for details.

What can be vacuumed?

- Locking a buffer for cleaning has special requirements (LockBufferForCleanup())
- It is safe for multiple backends to hold a reference to a buffer, and to do so for a long time
- This is not true when attempting to clean a page.
- Except for freezing, we will give up on trying to clean a page if any other backend is referencing it

Every time a backend takes a reference to a buffer, in gets a "pin". See src/backend/storage/buffer/README.

Long-running transactions prevent vacuuming from being effective

Freezing

- Transaction IDs (XIDs) and Multi-transaction IDs (MXIDs) eventually roll-over. Old ones must be frozen before this happens
- XIDs are created by transactions that modify data. MXIDs occur when more than one backend concurrently update/lock a row OR by SELECT FOR SHARE
- If either of these is in danger of rolling over, a special FREEZE vacuum must be run

See src/backend/access/heap/README.tuplock

Extremely high rates of update transactions, FOR SHARE LOCK, or concurrent FK checks can cause freeze problems

Heap Only Tuples

- Normal vacuum is quite expensive, but there are cases where we can avoid it
- If an update does not change any index values* and the new tuple will fit on the same page then we don't need to update indexes. We can update just the heap.
- Dead HOT tuples are optimistically removed when heap pages are read

* Indexed values means any column referenced anywhere in an index, including predicates and functions. See src/backend/access/heap/README.HOT.

Avoid referencing* heavily updated columns in indexes

* Indexed values means any column referenced anywhere in an index, including predicates and functions. See src/backend/access/heap/README.HOT.

VACUUM

- There are 4 major variations on vacuuming
- autovacuum is a built-in process that attempts to automatically vacuum anything that needs it
- VACUUM FULL completely rebuilds a table from scratch
- VACUUM is a regular, manually run vacuum
- VACUUM FREEZE is a manual vacuum that forcibly freezes everything it can

autovacuum

- Generally doesn't need to be tweaked in 9.4
- There is no way to control when it runs; do not attempt to do so with autovacuum_naptime
- If you do have slow periods (ie: weekends) it can help to run regular vacuum via cron
- Frequent manual vacuum of heavily updated tables is still a good idea
- Documented in sections 2.1.6 and 18.10

VACUUM FULL

- Since 9.0, completely rebuilds table and indexes from scratch, similar to cluster
- Takes an exclusive lock on table
- Because this is essentially a Create Table As Select + indexes, it's not really vacuuming anything
- See also https://github.com/reorg/pg_repack

VACUUM

- Can not be run in a transaction (or function)
- See also vacuumdb shell command
- For each table
 - Scan heap, remembering tuples to remove
 - Scan indexes, removing tuples
 - Remove tuples from heap
 - If ANALYZE option specified, do analyze.
- Update datfrozenxmin and datminmxid

If maintenance_work_mem isn't large enough, any vacuum (except full) must make multiple passes over indices

vacuum rel

- vacuum() -> vacuum_rel()
- Vacuums a single relation
- Does a bunch of mundane stuff, then calls either cluster_rel (for a VACUUM FULL) or lazy_vacuum_rel()
- Before returning, calls itself to vacuum the TOAST table (but not for autovac)

lazy_vacuum_rel

- vacuum() -> vacuum_rel() ->
 lazy_vacuum_rel()
- Decide if we need to freeze due to XID or MXID
- Scan the heap (and indexes): lazy_scan_heap()
- If it makes sense, lazy_truncate_heap()
- Clean up the free space map
- Update pg_class; log stats if needed

There was a bug in some versions where we updated relfrozenxid and relminmxid even if we hadn't scanned the whole table, potentially resulting in data loss.

lazy_scan_heap

• vacuum() -> vacuum_rel() -> lazy_vacuum_rel() -> lazy_scan_heap()

• For each block:

- If not freezing and more than 32 blocks are marked all visible, skip ahead
- If almost out of space, remove index tuples and known heap tuples
- Attempt cleanup lock; if not freezing and fail, skip block
- Prune page (same as HOT)
- Remember dead tuples (if indexes) or remove
- Update free space map and visibility map
- Update stats, last pass through indexes

lazy_cleanup_index

- vacuum() -> vacuum_rel() -> lazy_vacuum_rel() -> lazy_scan_heap() -> lazy_cleanup_index()
- Call index-specific cleanup method. These methods must scan the entire index, checking each index pointer against the list of remembered dead tuples

It is very difficult to reduce the size of a bloated index. Don't let bloat happen, and if it does, reindex.

lazy_vacuum_heap

```
    vacuum() -> vacuum_rel() ->
lazy_vacuum_rel() ->
lazy_scan_heap() ->
lazy_vacuum_heap()
```

- Removes marked-dead tuples from heap
- On each page, defragment page and record free space

vac_update_datfrozenxid

- Called by vacuum()
- Updates datfrozenxid and datminmxid
- If new values for either
 - Truncate Commit LOG files (pg_clog)
 - Update internal frozen XID and MXID info
 - MultiXact files (pg_multixact) are truncated during checkpoint

See ForceTransactionIdLimitUpdate() and vac_truncate_clog()

autovacuum

- Has two parts, launcher and workers
- The launcher prioritizes databases by
 - Most in need of XID freeze (usually none)
 - Most in need of MXID freeze (usually none)
 - Least recently autovacuumed, skipping any database vacuumed less than nap time ago.
- Multiple workers can work on the same database at once
- Workers will be canceled if they interfere with other backends
- Compare how many autovacuum workers are running against autovacuum_max_workers to see if autovaccum is running into problems.

autovacuum worker

- Get list of heap tables & materialized views that need vacuum or analyze
- Get list of TOAST tables that need vacuuming
- TEMP tables are ignored
- For each relation; attempt to get lock. Skip if unavailable (unless freeze)
- vac_update_datfrozenxid
- exit

autovacuum does not prioritize tables within a database autovacuum can become ineffective for high demand tables if too many large tables need vacuum at once

vacuum cost delay

- Well documented; please see section 18.4.4
- The critical idea is that once we hit (auto)vacuum_cost_limit we sleep for (auto)vacuum_cost_delay. Increasing limit speeds vacuum; increasing delay slows vacuum.
- Don't slow vacuum too much
- On many systems you should set page_dirty lower than page_miss

In closing...

- Long running transactions hurt vacuum
- High transaction rates, use of FOR SHARE LOCK and concurrent FK checks increase the need to FREEZE
- Indexes referencing heavily updated columns prevent HOT
- Make maintenance_work_mem large for vacuum
- It's very difficult to reduce the size of a bloated index
- autovacuum can only do so much

Questions?

jim.Nasby@BlueTreble.com